# Who moved my cheese? How storage systems deal with changes in their media

Gala Yadgar

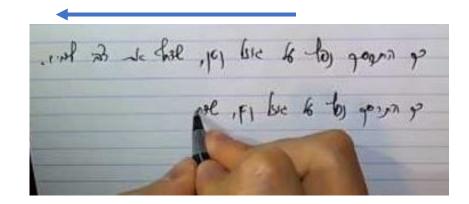


**TECHNION** | S The Henry and Marilyn Taub Faculty of Computer Science

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### "Why is Hebrew written backwards?"

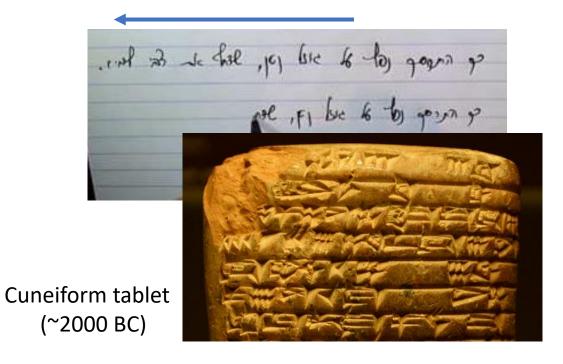
The quick brown fox jumps over the lazy dog.



## "Why is Hebrew written backwards?"

**Τhe quick brown fox μαρο ανεί τρε μα 20 αξα** ΝΚΑΝΑΙΝΑΒΑΙΨ ΝΙΜΑΝΝΑΜΑΓ ΚΑSΑS ΥΙΝΨΙSΓΑ Α ΕΙΨΑΝΑSΊΝΓΑΚΑ Ι SΥΙΑΥΑΝ ΝΙΒΑΓΑΝΚΨΙSΦΑΝΑ SΥΙΝΨΑΝΓΑΒΙΝ ΑΙΨ . GAHΨΑΝΑ SΥΙΝΨΑΝΓΑΒΙΝΑΙΨ. GAHΨΑΝΑ ΓΑΚΑΊS SIS YIAYAI . ΑΜΕΝΟΙΦΑ ΤΖΥΙSΨΑΤΕΙΑΑΑΑΤΑΑΓΑΕΤΑ ΣΑΨΑΤΑΓΚΑ ΥΛΟΚΙΤΕ SINNIM HANNEGAMNA

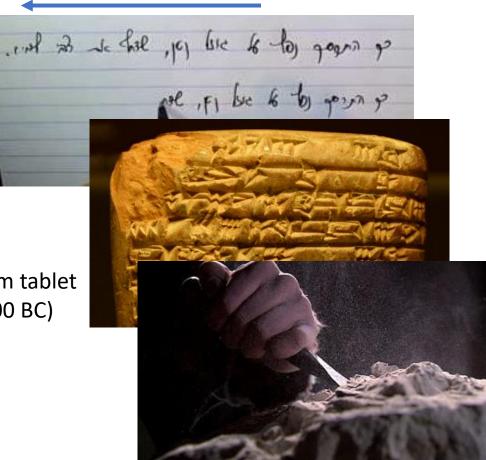
Codex Argenteus (~500 AC)



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## "Why is Hebrew written backwards?"

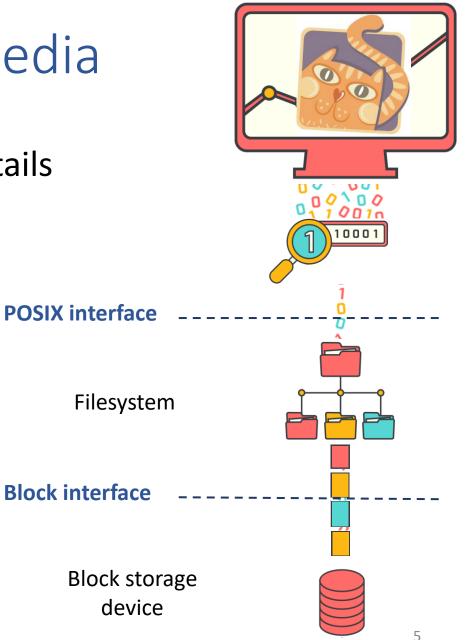




It's the storage media!

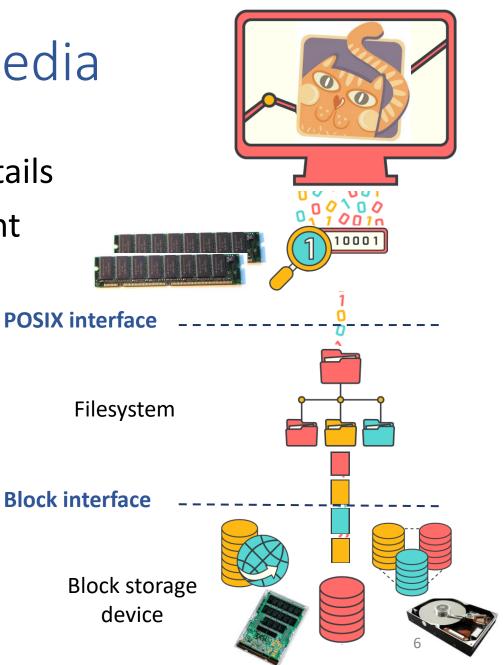
# Storage systems and their media

• Abstraction layers hide complexity and details



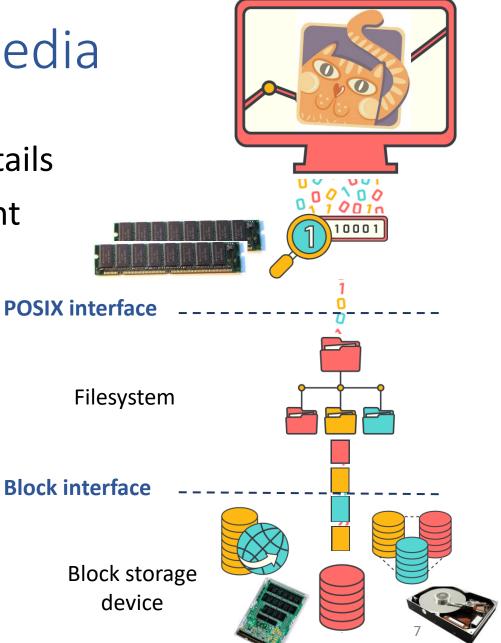
# Storage systems and their media

- Abstraction layers hide complexity and details
- Add the new media as a direct replacement
  - Backwards compatibility
  - Little development effort
  - Quick adoption



# Storage systems and their media

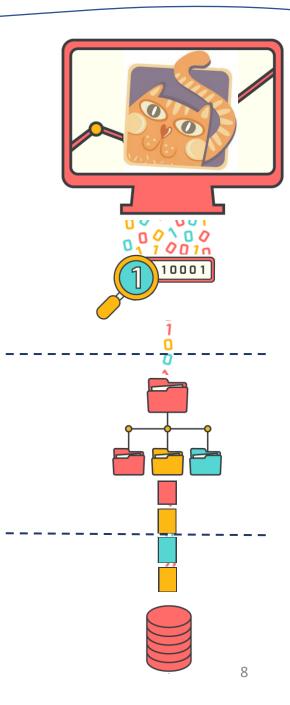
- Abstraction layers hide complexity and details
- Add the new media as a direct replacement
  - Backwards compatibility
  - Little development effort
  - Quick adoption
- It works! But...
  - Performance not as expected
  - Unpredictable user experience
- Notable examples from recent past
  - HDD  $\rightarrow$  RAID  $\rightarrow$  SSD  $\rightarrow$  NAS/cloud storage



# This talk

- How the storage-systems community addresses NVM
  Mon-exhaustive, non-prioritized list of examples\*
  Grossly over-simplified
  - Focus on insights and the adaptation process
- System model and characterization
- Unexpected bottlenecks
- Performance isolation / QoS

\* Figures taken from respective cited papers and author presentations



Non-volatile memory typically refers to storage in semiconductor memory chips, which store data in floating-gate memory cells consisting of floating-gate MOSFETs (metal-oxide-semiconductor field-effect transistors), including flash memory storage such as NAND flash and solid-state drives (SSD).

Other examples of non-volatile memory include read-only memory (ROM), EPROM (erasable programmable ROM) and EEPROM (electrically erasable programmable ROM), ferroelectric RAM, most types of computer data storage devices (e.g. disk storage, hard disk drives, optical discs, floppy disks, and magnetic tape), and early computer storage methods such as punched tape and cards.<sup>[1]</sup>



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  - Write-ahead logging
  - Metadata or index





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  - Write-ahead logging
  - Metadata or index
- For our purpose, byte-addresable and non-DRAM







 Intel Optane persistent memory

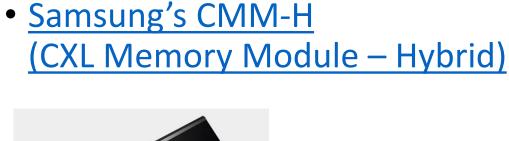


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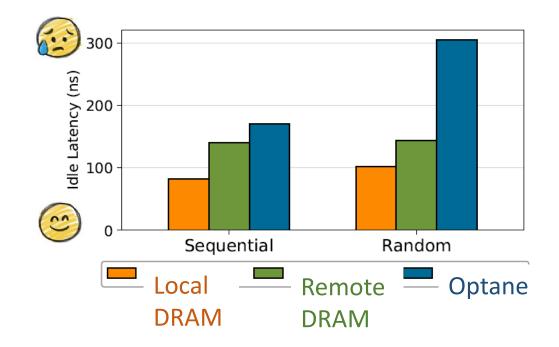






### Characterization (Optane)

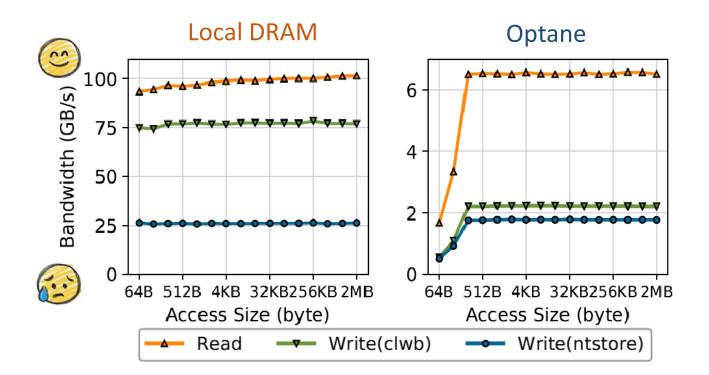
- Faster than HDD and SSD but slower than DRAM
- Does not behave like DRAM
  - Random ≠ Sequential



J. Izraelevitz et al. Basic Performance Measurements of the Intel Optane DC Persistent Memory Module, arXiv:1903.05714v3 2019

# Characterization (Optane)

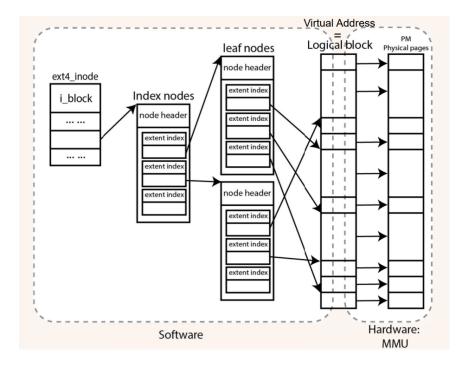
- Faster than HDD and SSD but slower than DRAM
- Does not behave like DRAM
  - Random ≠ Sequential
  - Reads ≠ Writes
  - Small ≠ Large
  - Much lower bandwidth
  - (Much more interference)



J. Izraelevitz et al. Basic Performance Measurements of the Intel Optane DC Persistent Memory Module, arXiv:1903.05714v3 2019

# Rethinking file-system indexing

- Bottleneck: tree-based inode index incurs high overhead with NVM
  - Up to 63% of the time spend on FS operations (e.g., file append)



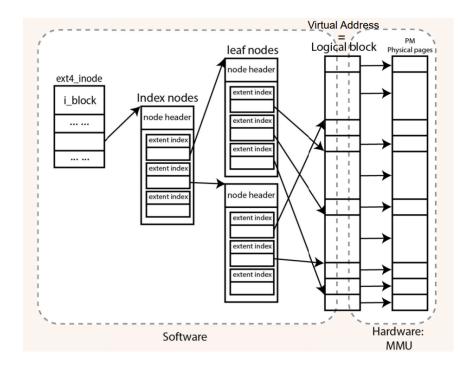
R. Li et al. <u>ctFS: Replacing File Indexing with Hardware Memory Translation through Contiguous File Allocation for</u> <u>Persistent Memory</u>, USENIX FAST 2022

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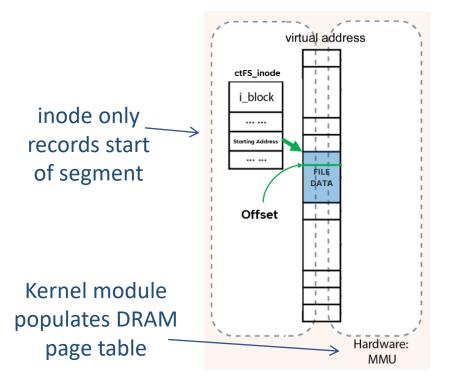
### • Insights:

- Use hardware-based translation
- Don't care about physical contiguity
- Persist small updates quickly



# Rethinking file-system indexing: ctFS

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# Rethinking file-system indexing: ctFS

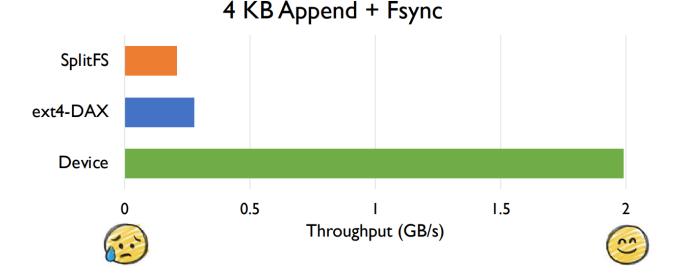
- Use hardware-based translation
- Don't care about physical contiguity
- New syscall: virtual address atomic pswap () Before pswap(A, B) After pswap(A, B) ctFS\_inode i block PMD PTE page PUD PUD PMD PTE page inode only ... ... records start arting Addr ... ... of segment FILE DATA Hierarchical Offset partitions Kernel module populates DRAM Hardware: page table MMU

Persist small updates quickly

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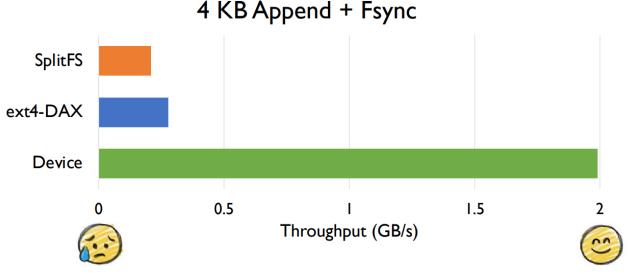
# Rethinking kernel-space vs. user-space

• Bottleneck: metadata handling in *kernel space* is inefficient



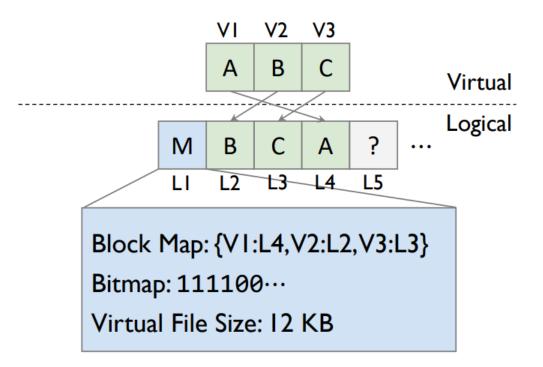
### Rethinking kernel-space vs. user-space

- Bottleneck: metadata handling in *kernel space* is inefficient
- Insights:
  - Some metadata can be safely handled in user space
  - Rethink data/metadata separation



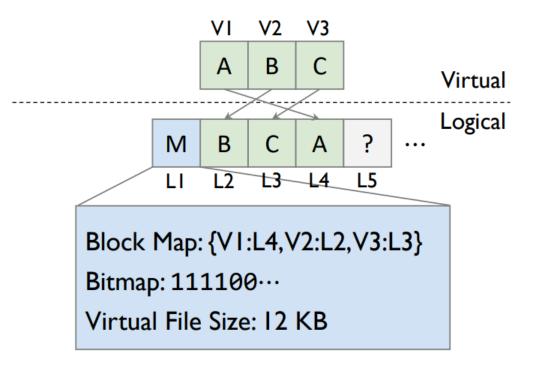
# Rethinking kernel-space vs. user-space: MadFS

- Some metadata can be safely handled in user space
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- Embedded metadata
  - "Virtual" file block map and size



# Rethinking kernel-space vs. user-space: MadFS

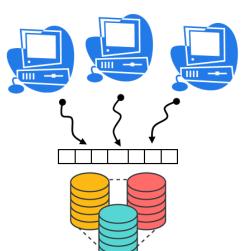
- Some metadata can be safely handled in user space
- Rethink data/metadata separation
- Embedded metadata
  - "Virtual" file block map and size
- Kernel-managed metadata
  - Logical-to-physical mapping
  - File permissions
  - Directory structures



# Rethinking shared caches

### • DRAM:

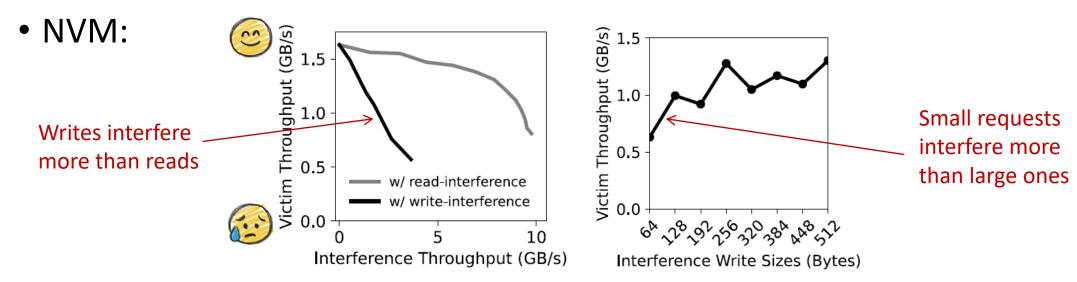
- More space allocation = higher hit rate = better performance
- Higher bandwidth = more usage = more interference
- Tenant A is too slow  $\rightarrow$  throttle noisy neighbor B with max bandwidth

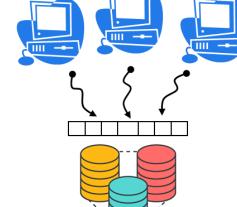


## Rethinking shared caches

### • DRAM:

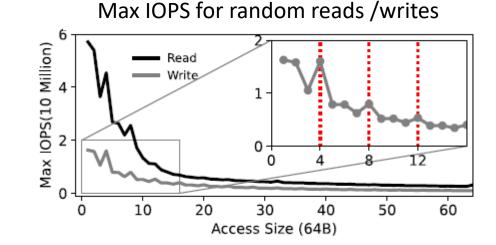
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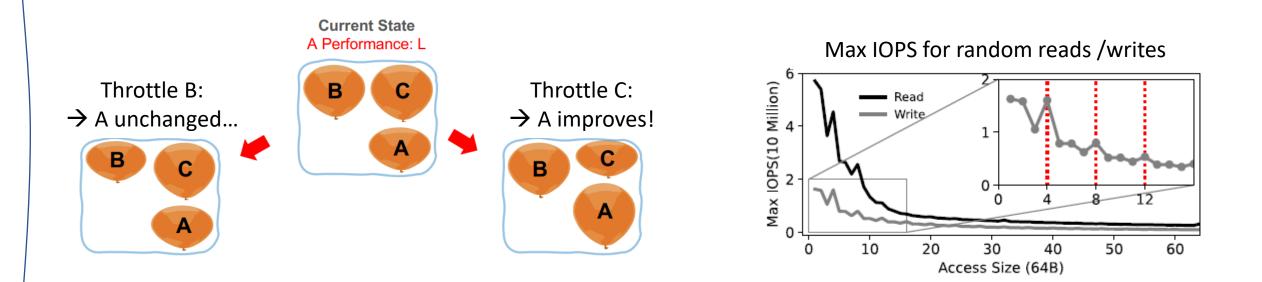
### Rethinking shared caches: NyxCache

• Profile max performance of different access types



## Rethinking shared caches: NyxCache

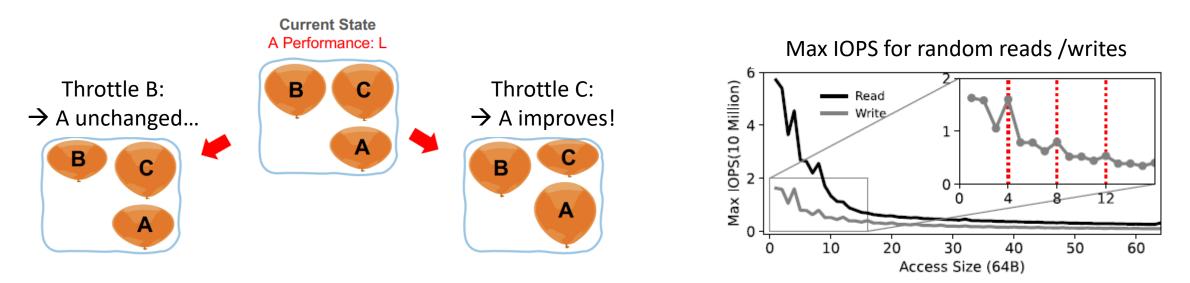
- Profile max performance of different access types
- Monitor individual tenants and their interaction

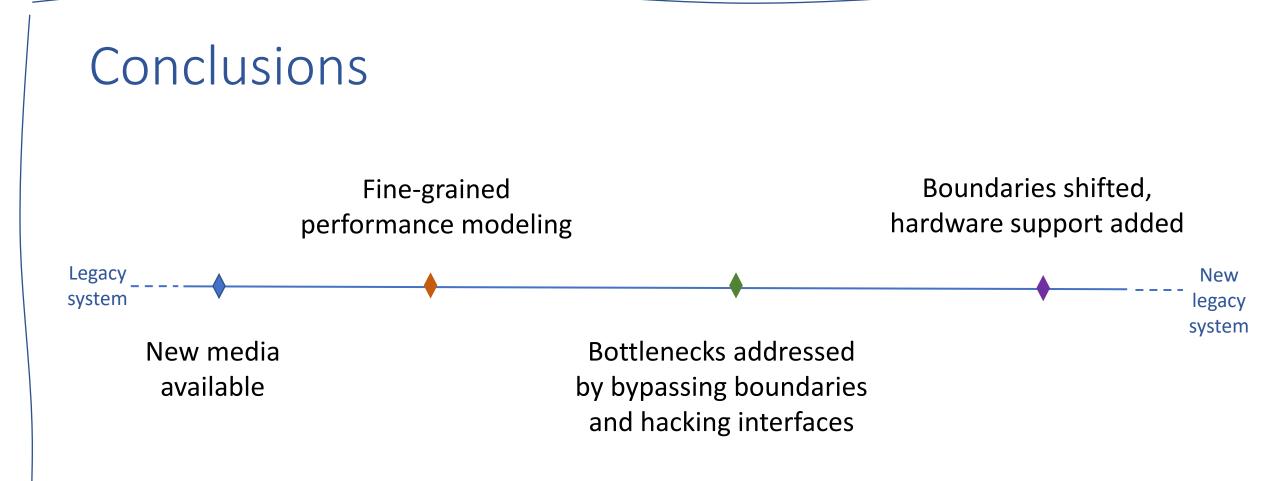


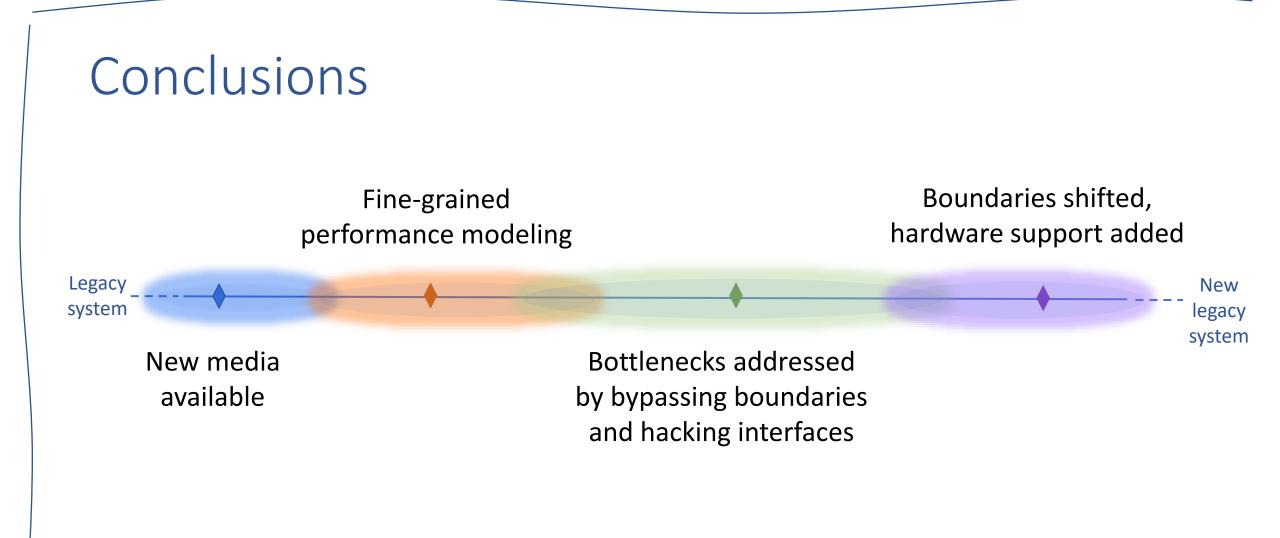
K. Wu at al. <u>NyxCache: Flexible and Efficient Multi-tenant Persistent Memory Caching</u> USENIX FAST 2022

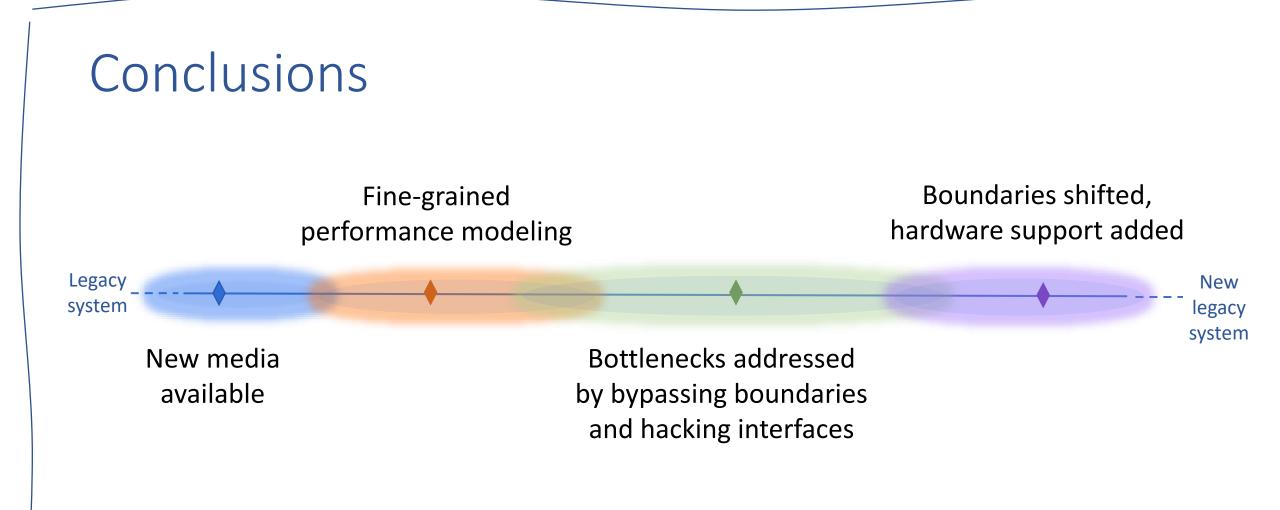
### Rethinking shared caches: NyxCache

- Profile max performance of different access types
- Monitor individual tenants and their interaction
- Library throttles tenant using most resources / causing most interference









#### • And what about our written language...?

